

# Fourier Based Fast Multipole Method for the Helmholtz Equation

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## Abstract

The fast multipole method (FMM) has had great success in reducing the computational time required to solve the boundary integral form of the Helmholtz equation. We present a formulation of the Helmholtz FMM using Fourier basis functions rather than spherical harmonics that accelerates some of the time-critical stages of the algorithm. With modifications to the transfer function in the precomputation stage of the FMM, the interpolation and antinterpolation operators become straightforward applications of fast Fourier transforms and the transfer operator remains diagonal. Using Fourier analysis, constructive algorithms are derived to a priori determine an integration quadrature for a given error tolerance. Sharp error bounds are derived and verified numerically. Various optimizations are considered to reduce the number of quadrature points and reduce the cost of computing the transfer function.

*Key words:* fast multipole method, fast Fourier transform, Fourier basis, interpolation, antinterpolation, Helmholtz, Maxwell, integral equations, boundary element method

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## 1 Introduction

2 Since the development of the fast multipole method (FMM) for the wave equa-  
3 tion in [1–5], the FMM has proven to be a very effective tool for solving scalar  
4 acoustic and vector electromagnetic problems. In this paper, we consider the  
5 application of the FMM to the scalar Helmholtz equations, although our re-  
6 sults can be immediately extended to the vector case as described in [6,7].

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7 The application of the boundary element method to solve the integral form of  
8 the Helmholtz equation results in a dense linear system which can be solved  
9 by iterative methods such as GMRES or BCGSTAB. These methods require  
10 computing dense matrix-vector products which, using a direct implementa-  
11 tion, are performed in  $\mathcal{O}(N^2)$  floating point operations. The FMM uses an  
12 approximation of the dense matrix to perform the product in  $\mathcal{O}(N \log N)$  or  
13  $\mathcal{O}(N \log^2 N)$  operations. This approximation is constructed from close pair in-  
14 teractions and far field approximations represented by spherical integrals that  
15 are accumulated and distributed through the domain via an octree.

16 There are a number of difficulties in implementing the FMM, each of which  
17 must be carefully considered and optimized to achieve the improved com-  
18 plexity. The most significant complication is that the quadrature sampling  
19 rate must increase with the size of the box in the octree, requiring interpo-  
20 lation and antepolation algorithms to transform the data between spherical  
21 quadratures of different levels of the tree. Local algorithms such as Lagrange  
22 interpolation and techniques which sparsify interpolant matrices are fast, but  
23 incur significant errors [8,7]. Spherical harmonic transforms are global inter-  
24 polation schemes and are exact but require fast versions for efficiency of the  
25 FMM. Many of these fast spherical transform algorithms are only approxi-  
26 mately, complicated to implement, and not always stable [9–11].

27 In this paper, we use a multipole expansion which allows the use of 2D fast  
28 Fourier transforms (FFT) in the spherical coordinate system  $(\phi, \theta)$ . The main  
29 advantages are two fold: i) high performance libraries are available for FFTs  
30 on practically all computer platforms, resulting in accurate, robust, and fast  
31 interpolation algorithms; ii) the resulting error analysis is simplified leading to  
32 sharp a priori error bounds on the calculation. One of the difficulties in using  
33 FFTs is that we are forced to use a uniform distribution of points along  $\phi$   
34 and  $\theta$  in the spherical quadrature. This leads to a much increased quadrature  
35 size for a given accuracy compared to the original spherical harmonics-based  
36 FMM. The reason is as follows. The multipole expansion in the high frequency  
37 regime is derived from:

$$\frac{e^{i\kappa|\mathbf{r}+\mathbf{r}_0|}}{|\mathbf{r} + \mathbf{r}_0|} = \int_{\phi=0}^{2\pi} \int_{\theta=0}^{\pi} e^{i\kappa\hat{\mathbf{s}}\cdot\mathbf{r}} T_{\ell,\mathbf{r}_0}(\hat{\mathbf{s}}) \sin(\theta) d\theta d\phi$$

38 where  $\hat{\mathbf{s}} = [\cos(\phi) \sin(\theta), \sin(\phi) \sin(\theta), \cos(\theta)]$  is the spherical unit vector. It  
39 is apparent that we are integrating along  $\theta$  a function which has period  $2\pi$ .  
40 However the bounds of the integral are 0 to  $\pi$ , over which interval the func-  
41 tion has a discontinuity in its derivative. This results in a slow decay of the  
42 Fourier spectrum (essentially  $1/\text{frequency}^2$ ) and consequently a large number  
43 of quadrature points along  $\theta$  are required.

44 We propose to use a variant of the scheme by J. Sarvas in [12] whereby the

45 integration is extended from 0 to  $2\pi$  and the integrand modified:

$$\frac{e^{i\kappa|\mathbf{r}+\mathbf{r}_0|}}{|\mathbf{r}+\mathbf{r}_0|} = \frac{1}{2} \int_{\phi=0}^{2\pi} \int_{\theta=0}^{2\pi} e^{i\kappa\hat{\mathbf{s}}\cdot\mathbf{r}} T_{\ell,\mathbf{r}_0}(\hat{\mathbf{s}}) |\sin(\theta)| d\theta d\phi$$

46 We will describe in more details how an efficient scheme can be derived from  
 47 this equation. The key property is that  $e^{i\kappa\hat{\mathbf{s}}\cdot\mathbf{r}}$  is approximately band-limited  
 48 in  $\theta$  and therefore it is possible to remove the high frequency components  
 49 of  $T_{\ell,\mathbf{r}_0}(\hat{\mathbf{s}}) |\sin(\theta)|$  without affecting the accuracy of the approximation. Using  
 50 this smooth transfer function, which is now band-limited in Fourier space, the  
 51 number of quadrature points can be reduced dramatically. We show that the  
 52 resulting number of quadrature points is reduced by about 40% compared to  
 53 the original spherical harmonics-based FMM. Consequently, we now have a  
 54 scheme which requires few quadrature points and enables the use of efficient  
 55 FFT routines. The approach in [12] is similar. However, rather than smooth-  
 56 ing  $T_{\ell,\mathbf{r}_0}(\hat{\mathbf{s}}) |\sin(\theta)|$  once during the precomputation phase as we will detail  
 57 shortly, Sarvas instead incorporates the  $|\sin(\theta)|$  factor during the run-time  
 58 phase of the FMM in the application of the transfer pass. This requires extra  
 59 interpolations/interpolations and a quadrature approximately twice the size  
 60 during the time-critical transfer pass of the algorithm. The method presented  
 61 in this paper retains a diagonal transfer pass operator and provides improved  
 62 control over the error.

63 We derive a new a priori error analysis which incorporates both effects from  
 64 truncation of the Gegenbauer series (a problem well analyzed [6]) and the  
 65 numerical quadrature. Our algorithm to predict the error is very sharp. This  
 66 allows choosing the minimal number of quadrature points to guarantee a pre-  
 67 scribed error. By comparison, the conventional approach can be shown to  
 68 result in conservative error bounds. That is, the number of quadrature points  
 69 is usually over-estimated and the error may be significantly below the target  
 70 accuracy. Although not considered in this paper, our error analysis approach  
 71 can also be applied to the spherical harmonics-based FMM to yield similarly  
 72 accurate error bounds. This has practical importance since it allows guaran-  
 73 teeing the error in the calculation while reducing the computational cost.

74 The paper is organized as follows. In section 2, we introduce the critical  
 75 parts of the classical FMM algorithm including the Gegenbauer series trunca-  
 76 tion (2.1), the spherical quadrature (2.2), and a short overview of interpola-  
 77 tion/interpolation strategies (2.3). In section 2.4, the asymptotic complexity  
 78 of the FMM is discussed. Section 3 details the Fourier basis approach. The  
 79 transfer function must be modified to lower the computational cost and obtain  
 80 a competitive scheme, as detailed in section 3.1. Section 3.2 applies an analysis  
 81 of the integration error to derive an algorithm that determines a quadrature  
 82 to yield an FMM with a prescribed error tolerance. The FFT based interpola-  
 83 tion and interpolation algorithms are described in section 3.3 and numerical  
 84 results are given in section 3.4. Table 1 lists the notations used in this paper.

Notation	Description
$\lambda$	wavelength
$\kappa$	wavenumber, $2\pi/\lambda$
$\theta$	polar angle
$\phi$	azimuthal angle
$\bar{a}$	complex conjugate of $a$
$\mathbf{x}$	vector in $\mathbb{R}^3$ , $\mathbf{x} =  \mathbf{x}  \hat{\mathbf{x}}$
$\mathbf{x} \cdot \mathbf{y}$	inner product, $\mathbf{x} \cdot \mathbf{y} =  \mathbf{x}   \mathbf{y}  \cos(\phi_{\mathbf{x},\mathbf{y}})$
$\mathbf{M}$	matrix in $\mathbb{R}^{n \times m}$ with elements $M_{ij}$
$S^2$	sphere, $\{\hat{\mathbf{s}} \in \mathbb{R}^3 :  \hat{\mathbf{s}}  = 1\}$
$j_n$	spherical Bessel function of the first kind
$y_n$	spherical Bessel function of the second kind
$h_n^{(1)}$	spherical Hankel function of the first kind
$P_n$	Legendre polynomial
$Y_n^m$	normalized spherical harmonic of degree $n$ , order $m$
$\mathcal{F}(m; f)$	$m$ th frequency of the Fourier transform of $f$

Table 1. Table of notations

## 85 2 The Multilevel Fast Multipole Method

86 The FMM reduces the computational complexity of the matrix-vector multi-  
87 plication

$$\sigma_i = \sum_{j \neq i} \frac{e^{i\kappa|\mathbf{x}_i - \mathbf{x}_j|}}{|\mathbf{x}_i - \mathbf{x}_j|} \psi_j = \sum_j M_{ij} \psi_j \quad (1)$$

88 for  $i, j = 1, \dots, N$  from  $\mathcal{O}(N^2)$  to  $\mathcal{O}(N \log^2 N)$ . This improvement is based on  
89 the Gegenbauer series

$$\frac{e^{i\kappa|\mathbf{r} + \mathbf{r}_0|}}{|\mathbf{r} + \mathbf{r}_0|} = i\kappa \sum_{n=0}^{\infty} (-1)^n (2n+1) h_n^{(1)}(\kappa|\mathbf{r}_0|) j_n(\kappa|\mathbf{r}|) P_n(\hat{\mathbf{r}} \cdot \hat{\mathbf{r}}_0) \quad (2)$$

90 The series converges absolutely and uniformly for  $|\mathbf{r}_0| \geq \frac{2}{\sqrt{3}} |\mathbf{r}|$  and has been  
91 studied extensively in [13,14].

92 Truncating the Gegenbauer series at  $\ell$  and using the identity

$$\int_{S^2} e^{i\kappa \hat{\mathbf{s}} \cdot \mathbf{r}} P_n(\hat{\mathbf{s}} \cdot \hat{\mathbf{r}}_0) dS(\hat{\mathbf{s}}) = 4\pi i^n j_n(\kappa |\mathbf{r}|) P_n(\hat{\mathbf{r}} \cdot \hat{\mathbf{r}}_0)$$

93 where the integral is over the unit sphere  $S^2$ , then

$$i\kappa \sum_{n=0}^{\ell} (-1)^n (2n+1) h_n^{(1)}(\kappa |\mathbf{r}_0|) j_n(\kappa |\mathbf{r}|) P_n(\hat{\mathbf{r}} \cdot \hat{\mathbf{r}}_0) = \int_{S^2} e^{i\kappa \hat{\mathbf{s}} \cdot \mathbf{r}} T_{\ell, \mathbf{r}_0}(\hat{\mathbf{s}}) dS(\hat{\mathbf{s}}) \quad (3)$$

94 where the transfer function,  $T_{\ell, \mathbf{r}_0}(\hat{\mathbf{s}})$ , is defined as

$$T_{\ell, \mathbf{r}_0}(\hat{\mathbf{s}}) = \frac{i\kappa}{4\pi} \sum_{n=0}^{\ell} i^n (2n+1) h_n^{(1)}(\kappa |\mathbf{r}_0|) P_n(\hat{\mathbf{s}} \cdot \hat{\mathbf{r}}_0). \quad (4)$$

95 Consider two disjoint clusters of points  $\{\mathbf{x}_i \mid i \in A\}$  and  $\{\mathbf{x}_i \mid i \in B\}$  with radii  
 96  $r_A \geq r_B \geq 0$  and centers  $\mathbf{c}_A$  and  $\mathbf{c}_B$  respectively. If  $|\mathbf{c}_A - \mathbf{c}_B| \geq \frac{2}{\sqrt{3}}(r_A + r_B)$ ,  
 97 then the matrix-vector product (1) is accelerated by using the approximation

$$\mathbf{M} \approx \begin{bmatrix} \mathbf{M}_{AA} & \widetilde{\mathbf{M}}_{AB} \\ \widetilde{\mathbf{M}}_{BA} & \mathbf{M}_{BB} \end{bmatrix}$$

98 where

$$\begin{aligned} \mathbf{M}_{AA} &= [M_{ij}]_{i \in A, j \in A} \\ \widetilde{\mathbf{M}}_{AB} &= \left[ \int_{S^2} e^{i\kappa \hat{\mathbf{s}} \cdot (\mathbf{c}_A - \mathbf{x}_i)} T_{\ell, \mathbf{c}_B - \mathbf{c}_A}(\hat{\mathbf{s}}) e^{i\kappa \hat{\mathbf{s}} \cdot (\mathbf{x}_j - \mathbf{c}_B)} dS(\hat{\mathbf{s}}) \right]_{i \in A, j \in B} \end{aligned}$$

99 For  $i \in A$ , this corresponds to computing

$$\begin{aligned} \sigma_i &= \int_{S^2} e^{i\kappa \hat{\mathbf{s}} \cdot (\mathbf{c}_A - \mathbf{x}_i)} T_{\ell, \mathbf{c}_B - \mathbf{c}_A}(\hat{\mathbf{s}}) \left[ \sum_{j \in B} e^{i\kappa \hat{\mathbf{s}} \cdot (\mathbf{x}_j - \mathbf{c}_B)} \psi_j \right] dS(\hat{\mathbf{s}}) \\ &\quad + \sum_{j \in A} \frac{e^{i\kappa |\mathbf{x}_i - \mathbf{x}_j|}}{|\mathbf{x}_i - \mathbf{x}_j|} \psi_j + \varepsilon_\ell \end{aligned}$$

100 where  $\varepsilon_\ell$  is the error introduced by the Gegenbauer series truncation.

101 The reduced computational complexity of the FMM is achieved by construct-  
 102 ing a tree of nodes, typically an octree, over the domain of the source and field  
 103 points. Let  $U_\alpha^l(\hat{\mathbf{s}})$  be the outgoing field for  $B_\alpha^l$ , the box  $\alpha$  of the tree in level  
 104  $l \in [0, L]$  with center  $\mathbf{c}_\alpha^l$ . The method is initialized by computing the outgoing  
 105 plane-wave expansions for each cluster contained in a leaf of the tree:

$$U_\alpha^L(\hat{\mathbf{s}}) = \sum_{i, \mathbf{x}_i \in B_\alpha^L} e^{i\kappa \hat{\mathbf{s}} \cdot (\mathbf{x}_i - \mathbf{c}_\alpha^L)} \psi_i$$

106 These outgoing expansions are then aggregated upward through the tree by  
 107 accumulating the product of the child cluster expansions with the diagonal  
 108 plane-wave translation function:

$$U_{\alpha}^l(\hat{\mathbf{s}}) = \sum_{\beta, B_{\beta}^{l+1} \subset B_{\alpha}^l} U_{\beta}^{l+1}(\hat{\mathbf{s}}) e^{i\kappa \hat{\mathbf{s}} \cdot (\mathbf{c}_{\beta}^{l+1} - \mathbf{c}_{\alpha}^l)}$$

109 Incoming plane-wave expansions,  $I_{\alpha}^l(\hat{\mathbf{s}})$  of box  $B_{\alpha}^l$ , are computed from the  
 110 outgoing by multiplication with the diagonal transfer function:

$$I_{\alpha}^l(\hat{\mathbf{s}}) = \sum_{\beta} U_{\beta}^l(\hat{\mathbf{s}}) T_{\ell, \mathbf{c}_{\beta}^l - \mathbf{c}_{\alpha}^l}(\hat{\mathbf{s}})$$

111 where the parent of  $B_{\beta}^l$  is a neighbor of the parent of  $B_{\alpha}^l$ , and  $B_{\beta}^l$  is not a  
 112 neighbor of  $B_{\alpha}^l$ . The incoming plane-waves are then disaggregated downward  
 113 through the tree to compute the local field  $D_{\alpha}^l(\hat{\mathbf{s}})$ :

$$D_{\alpha}^l(\hat{\mathbf{s}}) = D_{\beta}^{l-1}(\hat{\mathbf{s}}) e^{i\kappa \hat{\mathbf{s}} \cdot (\mathbf{c}_{\beta}^{l-1} - \mathbf{c}_{\alpha}^l)} + I_{\alpha}^l(\hat{\mathbf{s}})$$

114 where  $B_{\alpha}^l \subset B_{\beta}^{l-1}$ . At the finest level, the integration over the sphere is finally  
 115 performed and added to the near field contribution to determine the field value  
 116 at the  $N$  field points:

$$\sigma_i = \int_{S^2} D_{\alpha}^L(\hat{\mathbf{s}}) e^{i\kappa \hat{\mathbf{s}} \cdot (\mathbf{c}_{\alpha}^L - \mathbf{x}_i)} dS(\hat{\mathbf{s}}) + \sum_{j \neq i} \frac{e^{i\kappa |\mathbf{x}_i - \mathbf{x}_j|}}{|\mathbf{x}_i - \mathbf{x}_j|} \psi_j$$

117 where  $\mathbf{x}_i \in B_{\alpha}^L$  and  $\mathbf{x}_j \in B_{\beta}^L \cup B_{\alpha}^L$  where  $B_{\beta}^L$  is any neighbor of  $B_{\alpha}^L$ .

## 118 2.1 Truncation Parameter in the FMM

119 The truncation parameter  $\ell$  must be chosen so that the Gegenbauer series (2)  
 120 is converged to a desired accuracy. However, for  $n > x$ ,  $j_n(x)$  decreases super-  
 121 exponentially while  $h_n^{(1)}(x)$  diverges. The divergence of the Hankel function  
 122 causes the transfer function to oscillate wildly and become numerically un-  
 123 stable. Even though the expansion converges, roundoff errors will adversely  
 124 affect the accuracy if  $\ell$  is too large. Thus, while one must choose  $\ell > \kappa |\mathbf{r}|$   
 125 so that sufficient convergence is achieved for the plane-wave, it must also be  
 126 small enough to avoid the divergence of the transfer function.

127 The selection the truncation parameter  $\ell$  has been studied extensively and  
 128 and a number of procedures for selecting it have been proposed.

129 The empirical formula

$$\ell \approx \kappa |\mathbf{r}| + C(\varepsilon) \log(\pi + \kappa |\mathbf{r}|)$$

130 appears to have been first proposed by Rokhlin [2]. This was considered and  
 131 revised by Darve [14] using a detailed asymptotic analysis of the Gegenbauer  
 132 series.

133 The excess bandwidth formula (EBF) is derived from the convergence of the  
 134 plane-wave expansion and is presented in [6]. To determine an appropriate  
 135 truncature, the spectrum of a plane wave

$$e^{i\kappa\hat{\mathbf{s}}\cdot\mathbf{r}} = e^{i\kappa|\mathbf{r}|\cos(\phi_{\hat{\mathbf{s}},\mathbf{r}})} = \sum_{n=-\infty}^{\infty} i^n J_n(\kappa|\mathbf{r}|)e^{in\phi_{\hat{\mathbf{s}},\mathbf{r}}} \quad (5)$$

136 where  $\phi_{\hat{\mathbf{s}},\mathbf{r}}$  is the angle between  $\hat{\mathbf{s}}$  and  $\mathbf{r}$ , is used to estimate how many terms  
 137 in the series are needed before the error in the summation is exponentially  
 138 small. It can be shown that when  $n \rightarrow \infty$  and  $x \sim \mathcal{O}(n)$ ,

$$J_n(x) \sim \frac{e^{\sqrt{n^2-x^2}-n \cosh^{-1}(n/x)}}{\sqrt{2\pi(n^2-x^2)}}$$

139 which decays exponentially fast when  $n > x$ . Let  $n/x = 1 + \delta$  where  $\delta \ll 1$ .  
 140 Then  $\cosh^{-1}(n/x) \sim \sqrt{2\delta}$  and  $\sqrt{n^2-x^2} \sim x\sqrt{2\delta}$ . Thus, the above becomes

$$J_n(x) \sim \frac{e^{(x-n)\sqrt{2\delta}}}{2x\sqrt{\pi\delta}} = \frac{e^{-x\sqrt{2\delta^{3/2}}}}{2x\sqrt{\pi\delta}}$$

141 This expression is exponentially small when  $x\delta^{3/2} \gg 1$ , or  $\delta = Cx^{-2/3}$ , where  
 142  $C \gg 1$ . That is, when

$$\frac{n}{x} - 1 \approx Cx^{-2/3}$$

143 Therefore, the number of terms we need can be approximated as

$$\ell \approx \kappa|\mathbf{r}| + C(\kappa|\mathbf{r}|)^{1/3}$$

144 An empirically determined common choice is  $C = 1.8(d_0)^{2/3}$ , where  $d_0$  is the  
 145 desired number of digits of accuracy. The EBF is one of the most popular  
 146 choices to select the truncation parameter.

147 The desired number of digits of accuracy cannot always be achieved due to  
 148 the divergent nature of the transfer function. Thus, the EBF fails when high  
 149 accuracy is desired or the box size and box separation are small. Some modi-  
 150 fications in this regime are proposed in [15].

151 A direct numerical computation of the Gegenbauer truncation error  $\ell$  could  
 152 be computed or approximated

$$\varepsilon_G(\ell) = i\kappa \sum_{n=\ell+1}^{\infty} (-1)^n (2n+1) h_n^{(1)}(\kappa|\mathbf{r}_0|) j_n(\kappa|\mathbf{r}|) P_n(\hat{\mathbf{r}} \cdot \hat{\mathbf{r}}_0)$$

153 As Carayol and Collino showed in [13], an upper bound of this error for large  
 154 values of  $|\mathbf{r}|$  is obtained when  $P_n(\hat{\mathbf{r}} \cdot \hat{\mathbf{r}}_0) = P_n(\pm 1) = (\pm 1)^n$  so that

$$|\varepsilon_G| \lesssim \kappa \left| \sum_{n=\ell+1}^{\infty} (\mp 1)^n (2n+1) h_n^{(1)}(\kappa |\mathbf{r}_0|) j_n(\kappa |\mathbf{r}|) \right|$$

155 which they showed can be computed in closed form

$$\begin{aligned} &= \kappa \frac{\sqrt{|\mathbf{r}| |\mathbf{r}_0|}}{|\mathbf{r}_0| \pm |\mathbf{r}|} \frac{\pi}{2} \left| H_{\ell+\frac{3}{2}}(\kappa |\mathbf{r}_0|) J_{\ell+\frac{1}{2}}(\kappa |\mathbf{r}|) \pm H_{\ell+\frac{1}{2}}(\kappa |\mathbf{r}_0|) J_{\ell+\frac{3}{2}}(\kappa |\mathbf{r}|) \right| \\ &= \kappa^2 \frac{|\mathbf{r}| |\mathbf{r}_0|}{|\mathbf{r}_0| \pm |\mathbf{r}|} \left| h_{\ell+1}^{(1)}(\kappa |\mathbf{r}_0|) j_{\ell}(\kappa |\mathbf{r}|) \pm h_{\ell}^{(1)}(\kappa |\mathbf{r}_0|) j_{\ell+1}(\kappa |\mathbf{r}|) \right| \end{aligned}$$

156 This fails for small  $|\mathbf{r}|$  when the upper bound is instead given by the  $\hat{\mathbf{r}} \cdot \hat{\mathbf{r}}_0$   
 157 which causes the oscillation of  $P_n(\hat{\mathbf{r}} \cdot \hat{\mathbf{r}}_0)$  to compensate for the oscillation of  
 158  $(-1)^n h_n^{(1)}(\kappa |\mathbf{r}_0|) j_n(\kappa |\mathbf{r}|)$ .

159 Using the EBF as an initial guess for  $\ell$  and refining the choice using the above  
 160 closed form when  $|\mathbf{r}|$  is sufficiently large is a simple algorithm which yields an  
 161 accurate value for  $\ell$ .

162 Carayol and Collino in [16] and [13] present an in-depth analysis of the Jacobi-  
 163 Anger series and the Gegenbauer series. They find the asymptotic formula

$$\ell \approx v - \frac{1}{2} + \left(\frac{1}{2}\right)^{5/3} W^{2/3} \left( \frac{v}{4\varepsilon^6} \left(\frac{1+\alpha}{1-\alpha}\right)^{3/2} \right)$$

164 where  $v = \kappa |\mathbf{r}|$ ,  $\alpha = u/v$ ,  $u = \kappa |\mathbf{r}_0|$ , and  $W(x)$  is the Lambert function  
 165 defined as the solution to

$$W(x)e^{W(x)} = x \quad x > 0$$

166 This appears to be near optimal for large box sizes.

167 The errors introduced by this truncation have been investigated in other pa-  
 168 pers including [8,13,14].

## 169 2.2 Spherical Quadrature in the FMM

170 With the expansion formula

$$(2n+1)P_n(\hat{\mathbf{p}} \cdot \hat{\mathbf{q}}) = 4\pi \sum_{m=-n}^n \overline{Y_n^m(\hat{\mathbf{p}})} Y_n^m(\hat{\mathbf{q}})$$



171 the transfer function (4) can be expressed as

$$T_{\ell, \mathbf{r}_0}(\hat{\mathbf{s}}) = i\kappa \sum_{n=0}^{\ell} i^n h_n^{(1)}(\kappa |\mathbf{r}_0|) \sum_{m=-n}^n \overline{Y_n^m(\hat{\mathbf{r}}_0)} Y_n^m(\hat{\mathbf{s}}) = \sum_{n=0}^{\ell} \sum_{m=-n}^n t_n^m Y_n^m(\hat{\mathbf{s}}) \quad (6)$$

172 Similarly, the Jacobi-Anger series,

$$e^{i\kappa \hat{\mathbf{s}} \cdot \mathbf{r}} = \sum_{n=0}^{\infty} i^n (2n+1) j_n(\kappa |\mathbf{r}|) P_n(\hat{\mathbf{s}} \cdot \hat{\mathbf{r}}) \quad (7)$$

173 becomes the spherical harmonic series

$$e^{i\kappa \hat{\mathbf{s}} \cdot \mathbf{r}} = 4\pi \sum_{n=0}^{\infty} i^n j_n(\kappa |\mathbf{r}|) \sum_{m=-n}^n \overline{Y_n^m(\hat{\mathbf{r}})} Y_n^m(\hat{\mathbf{s}}) = \sum_{n=0}^{\infty} \sum_{m=-n}^n e_n^m Y_n^m(\hat{\mathbf{s}}) \quad (8)$$

174 The error analysis simplifies if we use a scheme which exactly integrates spheri-  
 175 cal harmonics,  $Y_l^m$ , up to some order. Below, we enumerate a number of choices  
 176 that have previously been studied.

- 177 (1) The simplest choice are sample points chosen uniformly in  $\phi$  and  $\theta$ . How-  
 178 ever, this choice does not accurately integrate the spherical harmonics  
 179 and requires approximately twice as many points as the Gauss-Legendre  
 180 quadrature below [7].
- 181 (2) The most common choice of sample points are uniform points for  $\phi$  and  
 182 Gauss-Legendre points for  $\theta$ . With  $N+1$  uniform points in the  $\phi$  direction  
 183 and  $\frac{N+1}{2}$  Gauss-Legendre points in the  $\theta$  direction, all  $Y_n^m$ ,  $-n \leq m \leq n$ ,  
 184  $0 \leq n \leq N$  are integrated exactly [7,8].
- 185 (3) McLaren in [17] developed optimal choices of samples for general func-  
 186 tions on  $S^2$  based on equally spaced points and derived from invariants  
 187 of finite groups of rotations. He also proposes a method for constructing  
 188 equally weighted integration formulas on sets of any desired number of  
 189 points by taking the union of icosahedral configurations.

### 190 2.3 Interpolation and Anterpolation in the FMM

191 The quadrature sampling rate depends on the spectral content of the diagonal-  
 192 ized translation operator,  $e^{i\kappa \hat{\mathbf{s}} \cdot \mathbf{r}}$ . These plane-waves contain more oscillatory  
 193 modes as we go up in the tree. Its coefficient in the spherical harmonic ex-  
 194 pansion  $e_n^m$  [equation (8)] decreases super-exponentially roughly for  $n \gtrsim \kappa |\mathbf{r}|$ .  
 195 Therefore, as we go up the tree in the aggregation step and  $|\mathbf{r}|$  becomes larger,  
 196 a larger quadrature is required to resolve these higher modes. These modes  
 197 must be resolved since they interact with the modes in the transfer function,  
 198 which do not significantly decay as  $\ell$  increases.

199 Similarly, as we go down the tree in the disaggregation step,  $|\mathbf{r}|$  becomes  
 200 smaller and the higher modes of the incoming field make vanishingly small  
 201 contributions to the integral as a consequence of Parseval’s theorem. Thus, as  
 202 the incoming field is disaggregated down the tree, a smaller quadrature can be  
 203 used to resolve it. This makes the integration faster and is actually required to  
 204 achieve an optimal asymptotic running time. See section 2.4 and appendix A.

205 There have been several approaches to performing the interpolation and an-  
 206 terpolation between levels in the FMM. Below, we enumerate a number of  
 207 choices that have previously been studied.

- 208 (1) General interpolation algorithms like Lagrange interpolation or B-splines  
 209 are fast and provide for simple error analysis. In [8] it is shown that the  
 210 error induced from Lagrange interpolation decreases exponentially as the  
 211 number of interpolation points is increased for a given function of finite  
 212 bandwidth. Thus, there is a trade-off between error and speed.
- 213 (2) For a set of quadrature points  $(\phi_k, \theta_k)$ ,  $k = 1, \dots, K$  with respective  
 214 weights  $\omega_k$  and corresponding function value  $f_k$ , a spherical harmonic  
 215 transform maps  $f_k$  to a new quadrature  $(\phi'_{k'}, \theta'_{k'})$ ,  $k' = 1, \dots, K'$  via the  
 216 linear transformation

$$f_{k'} = \sum_{m,l \leq K} Y_l^m(\phi'_{k'}, \theta'_{k'}) \sum_k \omega_k \overline{Y_l^m(\phi_k, \theta_k)} f_k = \sum_k A_{k'k} f_k \quad (9)$$

217 This transform has nice properties analogous to those of the Fourier trans-  
 218 form. A direct computation requires  $\mathcal{O}(KK')$  operations which would  
 219 result in an  $\mathcal{O}(N^2)$  FMM (see appendix A). Fast spherical transforms  
 220 (FST) have been developed in [9–11,18] and applied to the FMM in [19].  
 221 Using the FST reduces the interpolation and antinterpolation procedures  
 222 to  $\mathcal{O}(K \log^2 K)$ , which results in an  $\mathcal{O}(N \log^2 N)$  FMM. However, the  
 223 accuracy and stability of these algorithms remain in question.

- 224 (3) Approximations of the spherical transform have also been investigated in  
 225 [20,7]. The interpolation matrix  $A_{k'k}$  in (9) can be sparsified in a number  
 226 of ways to provide an interpolation/antinterpolation method that scales as  
 227  $\mathcal{O}(K)$  with controllable relative error.
- 228 (4) Many other interpolation schemes exist with varying running times and  
 229 errors. Rokhlin presents a fast polynomial interpolator based on the fast  
 230 multipole method in [21]. See also [22].

## 231 2.4 Asymptotic Complexity

232 In order to resolve a sufficient number of spherical harmonics, the number  
 233 of points in the  $\phi$  and  $\theta$  directions must be  $\mathcal{O}(\ell) = \mathcal{O}(\kappa a)$ , where  $a$  is the  
 234 side length of the box. Therefore, the total number of quadrature points is

235  $\mathcal{O}(\ell^2) = \mathcal{O}((\kappa a)^2)$ . If  $a_0$  is the side length of a box at the root of the octree  
 236 (level  $l = 0$ ), then the side length of a box at level  $l$  is  $a_l = 2^{-l}a_0$  and has  
 237  $\mathcal{O}((\kappa 2^{-l}a_0)^2)$  quadrature points.

238 With these parameters defined, the asymptotic complexity of the FMM can  
 239 be determined by carefully counting the number of operations required in each  
 240 step. This is done in detail in Appendix A and the results are discussed below.

241 In modeling a uniform distribution of point scatterers over a volumetric do-  
 242 main, we take  $\mathcal{O}(N) = \mathcal{O}((\kappa a_0)^3)$  and  $L \sim \log(N^{1/3})$  to achieve a total algo-  
 243 rithmic complexity of  $\mathcal{O}(N \log N)$  using fast global interpolation methods. By  
 244 using local methods, this can be reduced to  $\mathcal{O}(N)$ .

245 In modeling the scattering from the surface of an object using a uniform dis-  
 246 tribution of basis functions, we take  $\mathcal{O}(N) = \mathcal{O}((\kappa a_0)^2)$  and  $L \sim \log(N^{1/2})$  to  
 247 achieve a total algorithmic complexity of  $\mathcal{O}(N \log^2(N))$  using fast global inter-  
 248 polation methods. By using local methods, this can be reduced to  $\mathcal{O}(N \log N)$ .

249 It should be noted again that for a given, fixed  $\kappa a_0$  there is a minimum size  
 250 for the leaves of the tree. Below this critical size,  $h_n^{(1)}$  oscillates wildly causing  
 251 numerical instability in the transfer function. Therefore, when  $N$  is very large  
 252 and the number of levels is saturated, this analysis fails and the algorithm is  
 253 dominated by the  $\mathcal{O}(N^2)$  computation of the close field contribution, albeit  
 254 with orders of magnitude speedup over a direct method. In the case when  $\kappa a_0$   
 255 is too small, broadband FMMs have been developed as detailed in [23,24].  
 256 However, in many applications  $\kappa a_0$  is large enough to allow for all practi-  
 257 cal  $L$  and  $N$ . Furthermore, by keeping the number of points per wavelength  
 258 constant, the  $\mathcal{O}(N \log N)$  behavior can always be achieved.

### 259 3 Fourier Based Multilevel Fast Multipole Method

260 The Fourier based fast multipole method is based on the identity

$$\int_{S^2} e^{i\kappa \hat{\mathbf{s}} \cdot \mathbf{r}} T_{\ell, \mathbf{r}_0}(\hat{\mathbf{s}}) dS(\hat{\mathbf{s}}) = \int_0^{2\pi} \int_0^{2\pi} e^{i\kappa \hat{\mathbf{s}} \cdot \mathbf{r}} T_{\ell, \mathbf{r}_0}^s(\hat{\mathbf{s}}) d\phi d\theta \quad (10)$$

261 where  $\hat{\mathbf{s}} = [\cos(\phi) \sin(\theta), \sin(\phi) \sin(\theta), \cos(\theta)]$  and

$$T_{\ell, \mathbf{r}_0}^s(\hat{\mathbf{s}}) = \frac{1}{2} T_{\ell, \mathbf{r}_0}(\hat{\mathbf{s}}) |\sin(\theta)| \quad (11)$$

262 is the modified transfer function. Noting that the integrand is continuous and  
 263 periodic, this formulation of the problem suggests the use of the Fourier func-  
 264 tions  $\{e^{in\phi} e^{im\theta}\}$  which form an orthonormal basis of  $L^2([0, 2\pi] \times [0, 2\pi])$ . This

265 allows i) using two dimensional uniform quadratures; ii) fast Fourier trans-  
 266 forms in the interpolation and antepolation steps; iii) spectral arguments in  
 267 the error analysis. Of these advantages, the most important is that the FFT  
 268 interpolations and antepolations are fast and exact. Since there is no interpo-  
 269 lation error, only the finite quadrature and the truncation of the Gegenbauer  
 270 series introduce error to the final solution. Thus, the error analysis is simpli-  
 271 fied and we will determine in this paper precise bounds on the final error. In  
 272 fact, our error analysis is fairly general and can be extended to the classical  
 273 FMM with schemes that exactly integrate spherical harmonics (see direct and  
 274 fast global methods in section 2.3 and appendix A). The result is a fast, easy  
 275 to implement, and controllable version of the FMM, which we detail in the  
 276 following sections.

### 277 3.1 Computing the Modified Transfer Function

278 Select a uniform quadrature with points  $(\phi_i, \theta_j)$  defined by

$$\phi_i = 2\pi \frac{i}{N_\phi} \quad \theta_j = 2\pi \frac{j}{N_\theta}$$

279 Noting that the plane wave  $e^{i\kappa \hat{\mathbf{s}} \cdot \mathbf{r}}$  and  $T_{\ell, \mathbf{r}_0}^{\mathbf{s}}(\hat{\mathbf{s}})$  both have spherical symmetry,

$$f(\hat{\mathbf{s}}) \Big|_{\hat{\mathbf{s}}(\phi, \theta)} = f(\hat{\mathbf{s}}) \Big|_{\hat{\mathbf{s}}(\pi + \phi, 2\pi - \theta)},$$

280 the computational and memory cost are reduced by computing and storing  
 281 only half of the quadrature points.

282 Additionally, in an FMM with a single octree, there are 316 distinct transfer  
 283 vectors  $\mathbf{r}_0$  per level. By enforcing symmetries in the quadrature, the number of  
 284 modified transfer functions that must be precomputed is reduced. Specifically,  
 285 by requiring  $N_\theta$  to be a multiple of 2 and  $N_\phi$  to be a multiple of 4, we enforce  
 286 reflection symmetries in the  $z = 0$ ,  $x = 0$ ,  $y = 0$ ,  $x = y$ , and  $x = -y$   
 287 planes. This reduces the number of modified transfer functions that need to  
 288 be precomputed from 316 per level to 34 – saving a factor of 9.3 in memory  
 289 and costing a negligible permutation of the values of a computed modified  
 290 transfer function. See Fig. 1.

291 A key step to constructing a fast algorithm is to remove the high frequencies  
 292 in  $T_{\ell, \mathbf{r}_0}^{\mathbf{s}}(\hat{\mathbf{s}})$  whose contribution to the final result is negligible. This reduces  
 293 the number of needed quadrature points considerably. If  $T_{\ell, \mathbf{r}_0}^{\mathbf{s}}(\hat{\mathbf{s}})$  were simply  
 294 sampled, significant aliasing would occur unless we used an unreasonably large

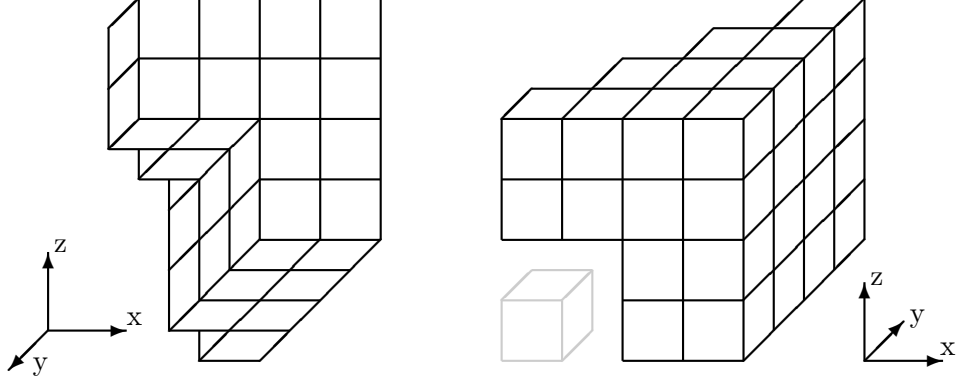


Fig. 1. The center of each box represents one transfer vector  $\mathbf{r}_0$  which must be computed. The pictures on the left and right panels represent the same set of boxes viewed under two different angles. Due to the symmetries of the quadrature, we need only compute transfer vectors with  $x, y, z \geq 0$  and  $x \geq y$ . We therefore end up with essentially half of an octant. Specifically, 34 transfer vectors are required; they can be reflected into any of the 316 needed.

295 quadrature. This is due to the slow decay of the Fourier series of  $|\sin(\theta)|$ ,

$$\mathcal{F}(m; |\sin(\theta)|) = \frac{(-1)^m + 1}{\pi(1 - m^2)} = \begin{cases} \frac{2}{\pi} \frac{1}{1 - m^2} & \text{if } m \text{ even} \\ 0 & \text{if } m \text{ odd} \end{cases}$$

296 Since the spectrum of the plane-wave function in equation (5) decays very  
 297 rapidly for  $n \gtrsim \kappa |\mathbf{r}|$ , the high frequencies in  $\theta$  of  $T_{\ell, \mathbf{r}_0}^{\mathbf{s}}(\hat{\mathbf{s}})$  do not contribute  
 298 to the final integral as a result of Parseval's theorem. By removing these  
 299 frequencies from the modified transfer function, a smaller quadrature can be  
 300 used without affecting the final result.

301 Suppose we have chosen a quadrature characterized by  $(N_\theta, N_\phi)$ . With this  
 302 quadrature we are able to exactly resolve the frequencies in  $e^{i\kappa \hat{\mathbf{s}} \cdot \mathbf{r}}$  between  
 303  $-N_\theta/2 + 1$  and  $N_\theta/2 - 1$  to the integral in equation (10). Consequently, we  
 304 need to exactly calculate a band limited approximation of  $T_{\ell, \mathbf{r}_0}^{\mathbf{s}}$ , called  $T_{\ell, \mathbf{r}_0}^{\mathbf{s}, L}$ ,  
 305 such that:

$$\mathcal{F}(m; T_{\ell, \mathbf{r}_0}^{\mathbf{s}, L}) = \begin{cases} \mathcal{F}(m; T_{\ell, \mathbf{r}_0}^{\mathbf{s}}), & \text{if } -N_\theta/2 + 1 \leq m \leq N_\theta/2 - 1 \\ 0, & \text{otherwise} \end{cases}$$

306 Since  $T_{\ell, \mathbf{r}_0}$  is bandlimited in  $\theta$  with bandwidth  $2\ell + 1$ , only the frequencies  
 307  $|m| \leq N_\theta/2 - 1 + \ell$  of  $|\sin(\theta)|$  contribute to the  $N_\theta - 1$  frequencies of  $T_{\ell, \mathbf{r}_0}^{\mathbf{s}, L}$ .  
 308 Therefore, the low-pass modified transfer function  $T_{\ell, \mathbf{r}_0}^{\mathbf{s}, L}$  can be computed using  
 309 the following pseudo-code:

```

1 for  $\phi_i$ ,  $0 \leq i < N_\phi/2$ , do
2    $T_k \leftarrow \frac{1}{2} T_{\ell, \mathbf{r}_0}(\phi_i, \frac{2\pi k}{2\ell+1})$ ,  $k = 0, \dots, 2\ell$ ;
3    $\tilde{T}_m \leftarrow \mathcal{F}(m, T)$ ;
310 4    $\tilde{s}_m \leftarrow \mathcal{F}(|m| \leq N_\theta/2 - 1 + \ell; |\sin(\theta)|)$ ;
5    $\tilde{T}_n^{\text{s},L} \leftarrow \tilde{s} \otimes \tilde{T}$  convolution of Fourier series;
6    $\tilde{T}_n^{\text{s},L} \leftarrow$  truncate to frequencies  $|n| \leq N_\theta/2 - 1$ ;
7    $T^{\text{s},L}(\theta_j, \phi_i) \leftarrow$  inverse transform of  $\tilde{T}_n^{\text{s},L}$ ;

```

311 This algorithm yields the low-pass modified transfer function at  $(\phi_i, \theta_j)$ ,  $0 \leq$   
312  $i < N_\phi/2$ ,  $0 \leq j < N_\theta$  which can be unwrapped to the remaining points  
313 by using the spherical symmetry  $(\phi_i, \theta_j) = (\phi_{N_\phi/2+i}, \theta_{N_\theta-j})$ . Note that this  
314 calculation can also be performed in the real space. It is equivalent to making  
315 a Fourier interpolation of  $T_k$  from  $2\ell+1$  points to  $N_\theta+2\ell-1$  points, multiplying  
316 by a low-pass  $|\sin(\theta)|$ , and performing a Fourier anteprolation back to  $N_\theta$   
317 points, as shown in Figure 2.

318 Because sampling the transfer function at a single point is an  $\mathcal{O}(\ell)$  operation,  
319 the algorithm as presented is  $\mathcal{O}(\ell^3)$ . The computation of the transfer function  
320 at all sample points can be accelerated to  $\mathcal{O}(\ell^2)$  as in [25] by taking advantage  
321 of its symmetry about the  $\hat{\mathbf{r}}_0$  axis and using interpolation algorithms, but at  
322 the cost of introducing additional error.

### 323 3.2 Choice of Quadrature

324 The quadrature parameters can be constructively computed by determining  
325 the maximum error they incur. The error in computing the integral with a  
326 finite uniform quadrature is

$$|\varepsilon_I| = \left| \int_0^{2\pi} \int_0^{2\pi} e^{i\kappa \hat{\mathbf{s}} \cdot \mathbf{r}} T_{\ell, \mathbf{r}_0}^{\text{s}}(\hat{\mathbf{s}}) d\phi d\theta - \sum_{n=1}^{N_\theta} \sum_{m=1}^{N_\phi} \omega_{n,m} e^{i\kappa \hat{\mathbf{s}}_{n,m} \cdot \mathbf{r}} T_{\ell, \mathbf{r}_0}^{\text{s},L}(\hat{\mathbf{s}}_{n,m}) \right|$$

where  $\hat{\mathbf{s}}_{n,m} = [\cos(\phi_m) \sin(\theta_n), \sin(\phi_m) \sin(\theta_n), \cos(\theta_n)]$  and  $T_{\ell, \mathbf{r}_0}^{\text{s},L}(\hat{\mathbf{s}}_{n,m})$  is the  
low-pass modified transfer function described in Section 3.1. This can be fur-  
327 ther expanded as:

$$\begin{aligned}
&= \left| \int_0^{2\pi} \int_0^{2\pi} \left[ E_{\kappa \mathbf{r}}^L(\hat{\mathbf{s}}_{n,m}) + E_{\kappa \mathbf{r}}^H(\hat{\mathbf{s}}_{n,m}) \right] \left[ T_{\ell, \mathbf{r}_0}^{\text{s},L}(\hat{\mathbf{s}}) + T_{\ell, \mathbf{r}_0}^{\text{s},H}(\hat{\mathbf{s}}) \right] d\phi d\theta \right. \\
&\quad \left. - \sum_{n=1}^{N_\theta} \sum_{m=1}^{N_\phi} \omega_{n,m} \left[ E_{\kappa \mathbf{r}}^L(\hat{\mathbf{s}}_{n,m}) + E_{\kappa \mathbf{r}}^H(\hat{\mathbf{s}}_{n,m}) \right] T_{\ell, \mathbf{r}_0}^{\text{s},L}(\hat{\mathbf{s}}_{n,m}) \right|
\end{aligned}$$

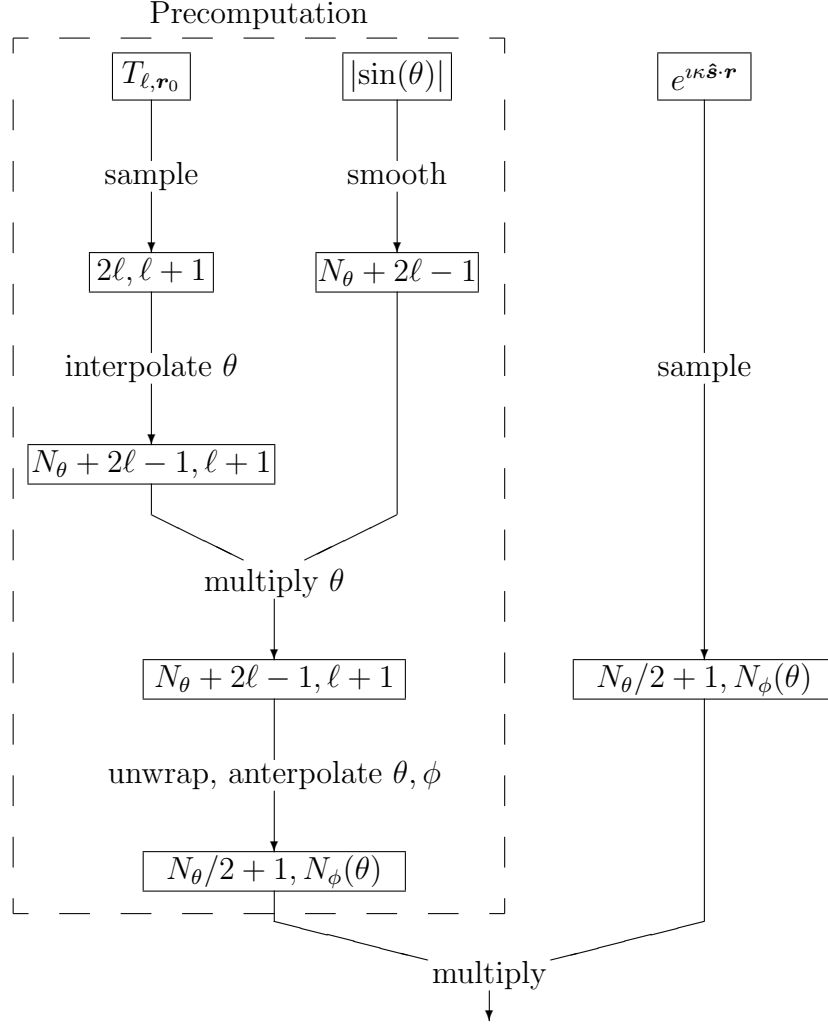


Fig. 2. Procedure for precomputing the low-pass modified transfer function and its application to an outgoing field. The boxed numbers (e.g.,  $2\ell, \ell + 1$ ) give the numbers of quadrature points for  $\theta$  and  $\phi$  ( $N_{\theta}$  and  $N_{\phi}$ ) at each stage.

where  $E_{\kappa \mathbf{r}}^L$  consists of the low frequencies of  $e^{i\kappa \hat{\mathbf{s}} \cdot \mathbf{r}}$  which are resolved by the quadrature and  $E_{\kappa \mathbf{r}}^H$  consists of the high frequencies that are aliased by the quadrature. Similarly,  $T_{\ell, \mathbf{r}_0}^{\mathbf{s}, H}$  are the high frequencies of  $T_{\ell, \mathbf{r}_0}^{\mathbf{s}}$  which we removed in section 3.1. Since  $E_{\kappa \mathbf{r}}^L T_{\ell, \mathbf{r}_0}^{\mathbf{s}, L}$  is integrated exactly by a uniform quadrature, and  $E_{\kappa \mathbf{r}}^H T_{\ell, \mathbf{r}_0}^{\mathbf{s}, L}$  does not contribute to the integral by Parseval's theorem, we get

$$\begin{aligned}
 &= \left| \int_0^{2\pi} \int_0^{2\pi} E_{\kappa \mathbf{r}}^H(\hat{\mathbf{s}}) T_{\ell, \mathbf{r}_0}^{\mathbf{s}, H}(\hat{\mathbf{s}}) d\phi d\theta - \sum_{n=1}^{N_{\theta}} \sum_{m=1}^{N_{\phi}} \omega_{n,m} E_{\kappa \mathbf{r}}^H(\hat{\mathbf{s}}_{n,m}) T_{\ell, \mathbf{r}_0}^{\mathbf{s}, L}(\hat{\mathbf{s}}_{n,m}) \right| \\
 &= \left| \int_0^{2\pi} \int_0^{2\pi} E_{\kappa \mathbf{r}}^H(\hat{\mathbf{s}}) T_{\ell, \mathbf{r}_0}^{\mathbf{s}, H}(\hat{\mathbf{s}}) - E_{\kappa \mathbf{r}}^{AH}(\hat{\mathbf{s}}) T_{\ell, \mathbf{r}_0}^{\mathbf{s}, L}(\hat{\mathbf{s}}) d\phi d\theta \right|
 \end{aligned}$$

329 where we have denoted the aliased high frequencies of  $e^{i\kappa\hat{\mathbf{s}}\cdot\mathbf{r}}$  as  $E_{\kappa\mathbf{r}}^{AH}(\hat{\mathbf{s}})$ ,

$$= \left| \int_0^{2\pi} \int_0^{2\pi} \left( E_{\kappa\mathbf{r}}^H(\hat{\mathbf{s}}) - E_{\kappa\mathbf{r}}^{AH}(\hat{\mathbf{s}}) \right) T_{\ell, \mathbf{r}_0}^s(\hat{\mathbf{s}}) d\phi d\theta \right|$$

330 In Fourier space, this becomes

$$= 4\pi^2 \left| \sum_{n=-\ell}^{\ell} \sum_{m=-\infty}^{\infty} \left( \tilde{E}_{\kappa\mathbf{r}}^H(n, m) - \tilde{E}_{\kappa\mathbf{r}}^{AH}(n, m) \right) \tilde{T}_{\ell, \mathbf{r}_0}^s(-n, -m) \right|$$

Choosing  $N_\theta$  using this expression for the error with a representative  $\mathbf{r}$  and  $\mathbf{r}_0$  leads to unpredictable cancellation effects and may result in a poor choice.

331 Instead, we apply the triangle inequality,

$$\leq 4\pi^2 \sum_{n=-\ell}^{\ell} \sum_{m=-\infty}^{\infty} \left| \tilde{E}_{\kappa\mathbf{r}}^H(n, m) - \tilde{E}_{\kappa\mathbf{r}}^{AH}(n, m) \right| \left| \tilde{T}_{\ell, \mathbf{r}_0}^s(-n, -m) \right|$$

332 This remains an accurate upper bound due to the fast decay of  $\tilde{E}$  for suffi-  
333 ciently large values of  $N_\theta$  and  $N_\phi$ . See Figure 3.

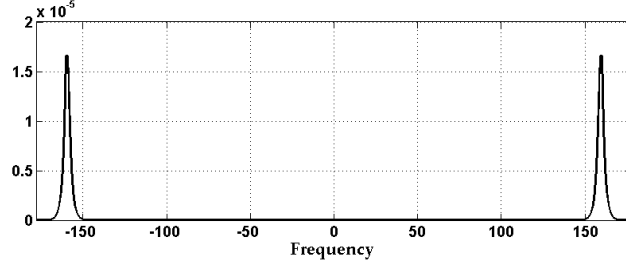


Fig. 3. The value of  $\left| \tilde{E}_{\kappa\mathbf{r}}^H(0, m) - \tilde{E}_{\kappa\mathbf{r}}^{AH}(0, m) \right|$  for  $\kappa |\mathbf{r}| = 0.8\sqrt{3} \cdot 100$  and  $N_\phi = 318$ .

### 334 3.2.1 Choosing $N_\theta$

335 The worst case for  $\varepsilon_I$  in terms of  $N_\theta$  occurs when  $\mathbf{r}$  and  $\mathbf{r}_0$  are aligned with the  
336  $z$ -axis. This causes all spectral information to be contained in the  $\theta$ -direction  
337 and makes  $\varepsilon_I$  a function of  $N_\theta$  only. It leads to

$$|\varepsilon_I| \leq 4\pi^2 \sum_{m=-\infty}^{\infty} \left| \tilde{E}_{\kappa\mathbf{r}}^H(0, m) - \tilde{E}_{\kappa\mathbf{r}}^{AH}(0, m) \right| \left| \tilde{T}_{\ell, \mathbf{r}_0}^s(0, -m) \right|$$

338 Using the plane wave spectrum (5) with  $\hat{\mathbf{r}} = \hat{\mathbf{z}}$ , this is approximately simplified  
339 to

$$|\varepsilon_I| \leq 4\pi^2 \sum_{m=-\infty}^{\infty} \left| J_{M(N_\theta, m)}(\kappa |\mathbf{r}|) \right| \left| \tilde{T}_{\ell, \mathbf{r}_0}^s(0, m) \right| \quad (12)$$



340 where

$$M(N_\theta, m) = \begin{cases} N_\theta - |m| & |m| \leq N_\theta/2 - 1 \\ |m| & |m| > N_\theta/2 - 1 \end{cases} \quad (13)$$

341 This is an approximation because  $\tilde{E}_{\kappa\mathbf{r}}^{AH}$  would in principle contribute an infi-  
 342 nite sum to equation (12) rather than the single term used. However, given  
 343 the exponential decay of the Jacobi-Anger series, the difference is negligible.  
 344 Equation (12) can be used to search for a value  $N_\theta$  via the algorithm sketched  
 345 below:

- 1 Choose  $N_\theta^n$  sufficiently larger than  $2\ell + 1$ ;
- 2  $T_k \leftarrow T_{\ell, |\mathbf{r}_0| \hat{\mathbf{z}}}^{s,L}(0, \frac{2\pi k}{N_\theta^n})$ ,  $k = 0, \dots, N_\theta^n - 1$ ;
- 3  $\tilde{T}_m \leftarrow |\mathcal{F}(m; T_k)|$ ;
- 4  $\tilde{E}_m \leftarrow |J_m(\kappa |\mathbf{r}|)|$ ;
- 5 **for**  $N_\theta$  from  $2\ell$  to  $N_\theta^n$  by 2 **do**
- 6      $\tilde{E}_m^* \leftarrow \tilde{E}_{M(N_\theta, m)}$ ;
- 7     **if**  $\tilde{E}^* \cdot \tilde{T} < \varepsilon/4\pi^2$  **then**
- 8         | return  $N_\theta$

346 Since  $N_\theta^n$  is typically only a small constant larger than  $2\ell + 1$ , the algorithm  
 347 as presented is dominated by the computation of the  $\mathcal{O}(\ell)$  modified transfer  
 348 function values and requires  $\mathcal{O}(\ell^2)$  operations. Important optimizations in-  
 349 clude using more advanced searching methods (such as bisection), applying  
 350 the symmetries  $\tilde{E}_m^* = \tilde{E}_{-m}^*$  and  $\tilde{T}_m = \tilde{T}_{-m}$ , and taking advantage of the very  
 351 fast decay of  $J_n$  to neglect very small terms in the dot product.

### 352 3.2.2 Choosing $N_\phi$

353 After determining an appropriate  $N_\theta$ , letting  $N_\phi$  be a function of  $\theta$  allows  
 354 reducing the number of quadrature points without affecting the error. The  
 355 worst case for the integration error occurs when  $\mathbf{r}$  and  $\mathbf{r}_0$  are in the  $xy$ -plane.  
 356 Without loss of generality, suppose  $\hat{\mathbf{r}} = \hat{\mathbf{x}}$ . Consider a constant  $\theta = \theta_j$  and  
 357 note that the plane wave can be expressed as

$$e^{i\kappa\hat{\mathbf{s}}\cdot\mathbf{r}} = \sum_{n=-\infty}^{\infty} i^n J_n(\kappa |\mathbf{r}| \sin(\theta_j)) e^{in\phi}$$

358 Since  $J_n(\kappa |\mathbf{r}| \sin(\theta_j))$  is exponentially small when  $n \gtrsim \kappa |\mathbf{r}| \sin(\theta_j)$ , the se-  
 359 ries can be truncated at  $N_\phi(\theta_j) \sim \kappa |\mathbf{r}| \sin(\theta_j)$  without incurring any ap-  
 360 preciable error. Estimates of  $N_\phi(\theta_j)$  can be developed by determining when  
 361  $J_n(\kappa |\mathbf{r}| \sin(\theta_j))$  becomes exponentially small, as in the computation of the  
 362 excess bandwidth formula (EBF) in [6]. However, we find that the EBF gen-  
 363 erated quadrature typically overestimates the sampling rate.

364 To accurately compute  $N_\phi(\theta_j)$  the same procedure as in Sec. 3.2.1 is applied  
 365 but with  $\mathbf{r}$  and  $\mathbf{r}_0$  in the  $xy$ -plane. This represents the worst case for the  
 366 integration error as a function of  $N_\phi$ . For a given  $\theta_j$ , we search for a  $N_\phi(\theta_j)$   
 367 such that

$$|\varepsilon_I| \leq 4\pi^2 \sum_{n=-\ell}^{\ell} \left| J_{M(N_\phi(\theta_j), n)}(\kappa |\mathbf{r}| \sin(\theta_j)) \right| \left| \tilde{T}_{\ell, \mathbf{r}_0}^s(n; \theta_j) \right| \quad (14)$$

368 is bounded by a prescribed error. This is accomplished via the following  
 369 sketched algorithm.

```

1 Set  $N_\phi$  at the poles:  $N_\phi(\theta_0) = N_\phi(\theta_{N_\theta/2}) = 1$ ;
2 Choose  $N_\phi^n$  sufficiently larger than  $2\ell + 1$ ;
3 for  $\theta_j, j = 1, \dots, N_\theta/2 - 1$  do
4    $T_k \leftarrow T_{\ell, |\mathbf{r}_0| \hat{\mathbf{x}}}^{\mathcal{S}, L}(\frac{2\pi k}{2\ell+1}, \theta_j), k = 0, \dots, 2\ell$ ;
5    $\tilde{T}_m \leftarrow |\mathcal{F}(m; T_k)|$ ;
6    $\tilde{E}_m \leftarrow |J_m(\kappa |\mathbf{r}| \sin(\theta_j))|$ ;
7   for  $N_\phi(\theta_j)$  from 2 to  $N_\phi^n$  by 2 do
8      $\tilde{E}_m^* \leftarrow \tilde{E}_{M(N_\phi(\theta_j), m)}$ ;
9     if  $\tilde{E}^* \cdot \tilde{T} < \varepsilon/4\pi^2$  then
10    | Save  $N_\phi(\theta_j)$ 
```

370 Since  $N_\phi^n$  is only a small constant larger than  $2\ell + 1$ , the algorithm as presented  
 371 is dominated by the computation of the modified transfer function and requires  
 372  $\mathcal{O}(\ell^3)$  operations. Optimizations similar to those presented in Sec. 3.2.1 can  
 373 be applied. Using the EBF as an initial guess in the search for  $N_\phi(\theta_j)$  further  
 374 improves the searching speed. Additionally, only half of the  $N_\phi(\theta_j)$ 's may be  
 375 computed due to symmetry about the  $z = 0$  plane.

376 We finally note that letting  $N_\phi$  be a function of  $\theta_j$  requires an additional step  
 377 in the computation of the modified transfer function. Section 3.1 computed  
 378 the transfer function on a  $N_\theta/2 + 1 \times N_\phi$  grid. With  $N_\phi \rightarrow N_\phi(\theta_j)$ , the data  
 379 computed for each  $\theta_j$  must be Fourier antepolated from length  $N_\phi$  to length  
 380  $N_\phi(\theta_j)$ .

### 381 3.2.3 Choosing $|\mathbf{r}|$ and $|\mathbf{r}_0|$

382 The previous algorithms require representative values of  $|\mathbf{r}|$  and  $|\mathbf{r}_0|$  for each  
 383 level of the tree. The worst-case transfer vectors,  $\mathbf{r}_0$ , are those with smallest  
 384 length. If  $a_l$  is the box size at level  $l$ , then  $|\mathbf{r}_0| = 2a_l$  is the smallest transfer  
 385 vector length in the common one buffer box case.

386 The worst case value of  $|\mathbf{r}|$  is the largest. For a box of size  $a_l$ ,  $|\mathbf{r}| \leq a_l\sqrt{3}$ .  
 387 However, using  $|\mathbf{r}| = a_l\sqrt{3}$  in the previous methods is often too conservative.

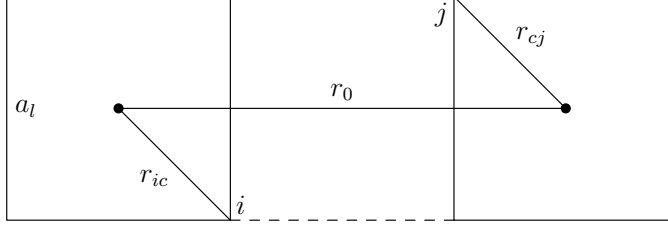


Fig. 4. The worst case  $\mathbf{r}$  and  $\mathbf{r}_0$ , projected from the 3D box. Here,  $|\mathbf{r}_0| = 2a_l$  and  $i$  and  $j$  are on the opposite corners of the box so that  $|\mathbf{r}| = |\mathbf{r}_{ic}| + |\mathbf{r}_{cj}| = a_l\sqrt{3}$ .

388 This case only occurs when two points are located in the exact corners of  
 389 the boxes – a rare case indeed. See Figure 4. Instead, we let  $|\mathbf{r}| = \alpha a_l\sqrt{3}$   
 390 for some  $\alpha \in [0, 1]$ . A high  $\alpha$  guarantees an upper bound on the error generated  
 391 by the quadrature, but the points which actually generate this error become  
 392 increasingly rare. A lower value of  $\alpha$  will yield a smaller quadrature, but more  
 393 points may fall outside the radius  $|\mathbf{r}|$  for which the upper bound on the error  
 394 is guaranteed.

### 395 3.2.4 Number of Quadrature Points

396 Recall from section 2.2 that the typical approach in the FMM is to use  $N + 1$   
 397 uniform points in the  $\phi$  direction and  $\frac{N+1}{2}$  Gauss-Legendre points in the  $\theta$   
 398 direction so that all  $Y_n^m$ ,  $-n \leq m \leq n$ ,  $0 \leq n \leq N$  are integrated exactly.  
 399 In [8], Chew et al. takes  $\frac{N+1}{2} = \ell + 1$ , which is an approximate choice based  
 400 on the rapid decay of the coefficients in the spherical harmonics expansion of  
 401 a plane wave. This results in approximately

$$M_g = 2(\ell + 1)^2 \approx 2\ell^2$$

402 quadrature points.

403 For a given Gegenbauer series truncation  $\ell$ , the total number of quadrature  
 404 points required in the Fourier based FMM is approximately

$$\begin{aligned} M_f &\approx \frac{N_\theta}{2} \frac{1}{\pi} \int_0^\pi N_\phi(\theta) d\theta \\ &\approx (\ell + C_1) \frac{1}{\pi} \int_0^\pi (2\ell + C_2(\theta)) \sin(\theta) d\theta \end{aligned}$$

405 where  $C_1, C_2 \geq 1$  are small integers dependent on  $\ell$ , numerically computed  
 406 from the methods in Sec. 3.2.1, 3.2.2. Keeping only the leading term in  $\ell$ :

$$M_f \approx \frac{4}{\pi} \ell^2 \approx 1.3 \ell^2$$

407 Thus, the method presented in this paper uses approximately 0.64 times the  
 408 number of quadrature points in the standard FMM. However, it is possible

409 that the same  $N_\phi$  optimization can be applied to the standard FMM for the  
 410 same reasons it was applied in section 3.2.2 to reduce the standard quadrature  
 411 to a comparable size.

### 412 3.3 Interpolation and Anterpolation

413 Most importantly, the Fourier based FMM directly uses FFTs in the inter-  
 414 polation and anterpolation steps. This makes the time critical upward pass  
 415 and downward pass especially fast and easy to implement while retaining the  
 416 exactness of global methods.

417 Characterize a quadrature by an array of length  $N_\theta/2 + 1$ ,

$$Q = [1, N_\phi(\theta_1), \dots, N_\phi(\theta_{N_\theta/2-1}), 1]$$

418 noting that  $N_\phi(\theta_j) = N_\phi(\theta_{N_\theta/2+j})$  and  $N_\phi(\theta_j) = N_\phi(\theta_{N_\theta/2-j})$ . The data  $F(\phi_i, \theta_j)$   
 419 sampled on a quadrature  $Q$  is transformed to a another quadrature  $Q'$  by per-  
 420 forming a sequence of Fourier interpolations and anterpolations. Let

$$\mathcal{N}_\phi = \max \left[ \max_{0 \leq j \leq N_\theta/2} N_\phi(\theta_j), \max_{0 \leq j \leq N'_\theta/2} N'_\phi(\theta_j) \right]$$

421 Then, the following steps, as illustrated in Figure 5, perform an exact inter-  
 422 polation/anterpolation using only FFTs.

- 423 (1) For each  $\theta_j$ ,  $0 \leq j \leq N_\theta/2$ , Fourier interpolate the data  $[F(\phi_{i=0, \dots, N_\phi(\theta_j)-1}, \theta_j)]$   
 424 from length  $N_\phi(\theta_j)$  to  $\mathcal{N}_\phi$ .
- 425 (2) For each  $\phi_i$ ,  $0 \leq i < \mathcal{N}_\phi/2$ , wrap the data to construct the periodic se-  
 426 quence from the rest of the line  $[F(\phi_i, \theta_{j=0, \dots, N_\theta/2}), F(\phi_{i+\mathcal{N}_\phi/2}, \theta_{j=N_\theta/2-1, \dots, 1})]$ .
- 427 (3) For each  $\phi_i$ ,  $0 \leq i < \mathcal{N}_\phi/2$ , Fourier interpolate the data  $[F(\phi_i, \theta_{j=0, \dots, N_\theta-1})]$   
 428 from length  $N_\theta$  to  $N'_\theta$ .
- 429 (4) For each  $\phi_i$ ,  $0 \leq i < \mathcal{N}_\phi/2$ , unwrap the data  $[F(\phi_i, \theta_{j=0, \dots, N'_\theta-1})]$  to con-  
 430 struct the sequences  $[F(\phi_i, \theta_{j=0, \dots, N'_\theta/2})]$  and  $[F(\phi_{i+\mathcal{N}_\phi/2}, \theta_{j=0, \dots, N'_\theta/2})]$ .
- 431 (5) For each  $\theta_j$ ,  $0 \leq j \leq N'_\theta/2$ , Fourier anterpolate the data  $[F(\phi_{i=0, \dots, \mathcal{N}_\phi(\theta_j)-1}, \theta_j)]$   
 432 from length  $\mathcal{N}_\phi$  to  $N'_\phi(\theta_j)$ .

### 433 3.4 Numerical Results

#### 434 3.4.1 Error

435 A direct computation was used to compute the optimal Gegenbauer truncation  
 436  $\ell$  and the methods described in section 3.2 were used to construct a quadrature  
 437 for use in computing the integral (10). For a given box size  $a$ , the quadrature

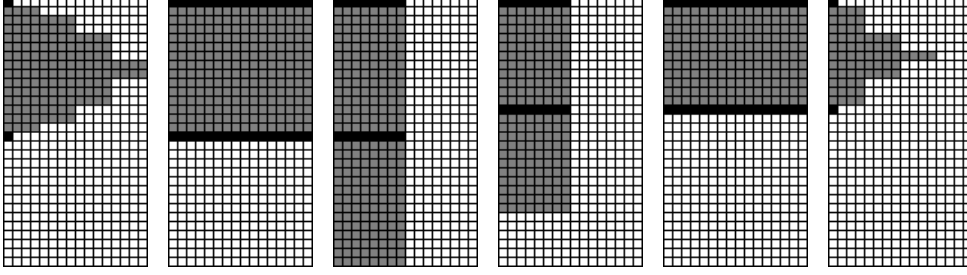


Fig. 5. The data profile at each step in an interpolation from a large quadrature  $Q$  with  $N_\theta = 30$  to a smaller quadrature  $Q'$  with  $N'_\theta = 24$ . The angle  $\phi$  is in the  $x$  direction while the angle  $\theta$  is in the  $y$  direction. The data corresponding to a pole has been darkened for clarity.

438 and truncation are constructed with  $|\mathbf{r}| = 0.8a\sqrt{3}$ ,  $|\mathbf{r}_0| = 2a$ , and target error  
 439  $\mathbf{eps}$ . The total measured error,  $\varepsilon$ , is defined as

$$\varepsilon = \frac{e^{i\kappa|\mathbf{r}+\mathbf{r}_0|}}{|\mathbf{r} + \mathbf{r}_0|} - \sum_{m=1}^{N_\theta} \sum_{n=1}^{N_\phi(\theta_m)} \omega_{n,m} e^{i\kappa\hat{\mathbf{s}}_{n,m}\cdot\mathbf{r}} T_{\ell,\mathbf{r}_0}^{\mathbf{s},L}(\hat{\mathbf{s}}_{n,m})$$

440 The total Gegenbauer truncation error,  $\varepsilon_G$ , is

$$\varepsilon_G = \frac{e^{i\kappa|\mathbf{r}+\mathbf{r}_0|}}{|\mathbf{r} + \mathbf{r}_0|} - i\kappa \sum_{n=0}^{\ell} (-1)^n (2n+1) h_n^{(1)}(\kappa|\mathbf{r}_0|) j_n(\kappa|\mathbf{r}|) P_n(\hat{\mathbf{r}} \cdot \hat{\mathbf{r}}_0)$$

441 The total integration error  $\varepsilon_I$  is

$$\varepsilon_I = \varepsilon - \varepsilon_G$$

442 In Figure 6, the plotted errors represent the maximum found over many direc-  
 443 tions  $\hat{\mathbf{r}}$  and magnitudes  $|\mathbf{r}| \leq 0.8s\sqrt{3}$ . As is evident, as the box size increases,  
 444 the target error  $\mathbf{eps}$  is accurately achieved. The increase in error for small box  
 445 sizes corresponds to the low frequency breakdown when the transfer function  
 446 has very large amplitude and roundoff errors become dominant. In this regime  
 447 the quadrature target error bound is also relaxed to improve efficiency - it is  
 448 inefficient to have a large quadrature that provides a small integration error  
 449 when the transfer function cannot provide comparable accuracy.

450 On the same plot we show  $\varepsilon_G^{EBF}$ , the Gegenbauer series error resulting from  
 451 choosing the truncation with the EBF from section 2.1. Clearly, the EBF  
 452 is overestimating  $\ell$ , which causes the Gegenbauer error to fall far below the  
 453 target error and will force the quadrature to be larger and less efficient.

454 The bottom plot shows the ratio of the number of points in the quadrature  
 455 presented in this paper to the number of quadrature points that would be  
 456 used in a typical spherical harmonics based FMM. Each of these quadratures  
 457 were computed for the same Gegenbauer truncation  $\ell$  chosen by the direct  
 458 calculation. The procedures presented in this paper result in a quadrature

459 which is substantially smaller than what would typically be used. Notably,  
 460 the analysis in Section 3.2.4 is supported.

461 Together, these results demonstrate that by choosing  $\ell$  and the quadrature  
 462 as presented in this paper, the error is better controlled and the quadrature  
 463 size at each level in the tree is reduced. Improved error control means that we  
 464 can provide a sharp bound of the total final error of the method and optimize  
 465 the running time of the method for that prescribed error. A reduction in the  
 466 quadrature size improves memory usage and suggests an improved running  
 467 time over similar algorithms.

### 468 3.4.2 Speed

469 As discussed in Section 3.3, the Fourier based FMM uses only FFTs in the  
 470 upward pass and downward pass to perform the interpolations and anterpo-  
 471 lations. FFTs make these steps easier to implement and very fast.

472 Figure 7 shows the recorded running times of the Fourier based FMM and the  
 473 direct matrix-vector product on a Intel Core 2 Quad CPU Q9450 2.66GHz  
 474 with 4GB of RAM. For  $N = 8.2 \cdot 10^6$  the points are uniformly distributed  
 475 in a cube with side length  $80\lambda$  and the wave number  $\kappa$  is scaled with  $N^{1/3}$ .  
 476 This provides a nearly constant density of points per wavelength as  $N$  is  
 477 varied. As expected, by choosing the correct number of levels the running  
 478 time is asymptotically  $\mathcal{O}(N \log N)$  as  $N$  is increased with a constant number  
 479 of points per wavelength. Note that the cross-over point is less than  $N = 4,000$ .  
 480 The code used to produce these results was not optimized for memory usage,  
 481 preventing results for  $N \gtrsim 10^6$  when  $L = 7$ .

## 482 4 Conclusion

483 We have proposed using the Fourier basis  $e^{ip\phi} e^{iq\theta}$  in the spherical variables  $\phi$   
 484 and  $\theta$  to represent the far field approximation in the FMM. By approximating  
 485 the Helmholtz kernel with

$$\frac{e^{i\kappa|\mathbf{r}+\mathbf{r}_0|}}{|\mathbf{r}+\mathbf{r}_0|} \approx \int_0^{2\pi} \int_0^{2\pi} e^{i\kappa\hat{\mathbf{s}}\cdot\mathbf{r}} T_{\ell,\mathbf{r}_0}^s(\hat{\mathbf{s}}) d\phi d\theta, \quad T_{\ell,\mathbf{r}_0}^s(\hat{\mathbf{s}}) = \frac{1}{2} T_{\ell,\mathbf{r}_0}(\hat{\mathbf{s}}) |\sin(\theta)|,$$

486 and using a uniform quadrature we can take advantage of very fast, exact, and  
 487 well-known FFT interpolation/anterpolation methods. By exploiting symme-  
 488 tries and a scheme to reduce the number of points in the  $\phi$  direction, the total  
 489 number of uniform quadrature points required is smaller than the number of  
 490 Gauss-Legendre quadrature points typically used with spherical harmonics.

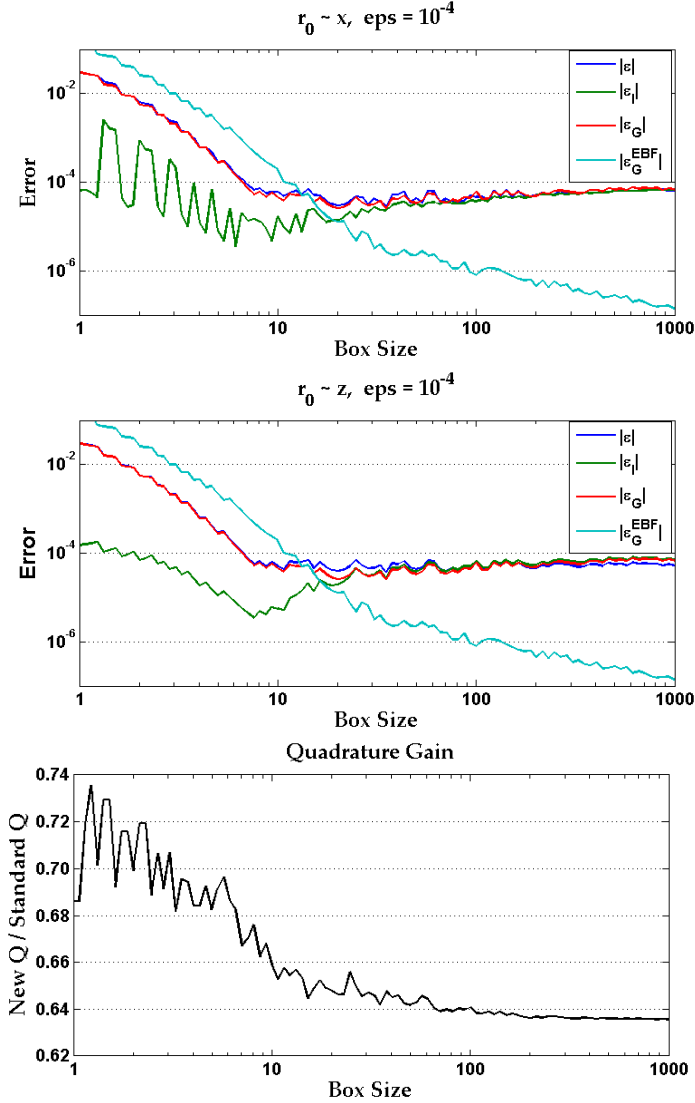


Fig. 6. Top plots: error of the FMM integral using a direct computation of  $\ell$  as described in section 2.1 and the choice of quadrature described in section 3.2. All the errors fall very close to the target error of  $10^{-4}$ . The standard EBF overestimates  $\ell$  and will result in a suboptimal quadrature. Bottom plot: ratio of the number of quadrature points required in the Fourier based FMM with what would be used in a typical spherical harmonics based FMM for the same  $\ell$ . The curve asymptotes close to  $2/\pi \approx 0.64$  as expected.

491 This is realized by correcting the transfer function  $T_{\ell, r_0}^s(\hat{\mathbf{s}})$  during the pre-  
 492 computation phase to remove high frequency terms which do not significantly  
 493 contribute to the final integration.

494 The Fourier based FMM approach has a number of advantages. Since the in-  
 495 terpolation and antinterpolation algorithms are exact, the error analysis is sim-  
 496 plified; we establish a sharp upper bound for the error. The key parameters are  
 497 the Gegenbauer truncation parameter  $\ell$  and the quadrature size, in particular

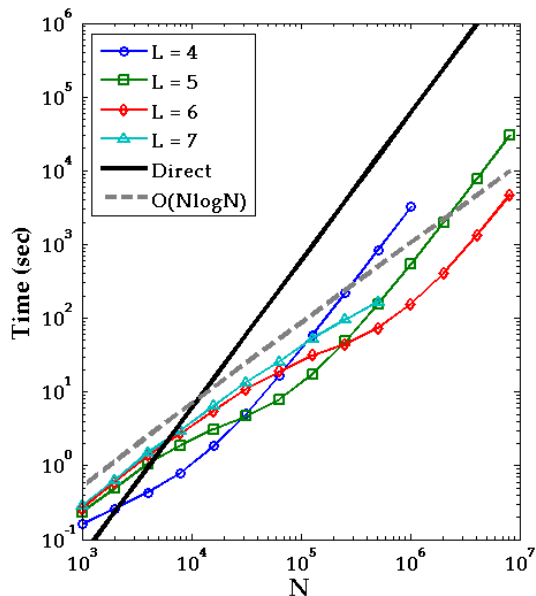


Fig. 7. Average running times of the Fourier based FMM for constant number of points per wavelength.

498 the bandwidth in the  $\theta$ -direction. The truncation error  $\varepsilon_\ell$  has been exten-  
 499 sively studied by other authors and is well understood. The integration error  
 500  $\varepsilon_I$  accounts for the low-pass approximation of the modified transfer function  
 501 and the aliasing of the plane waves. This error can be accounted for a priori  
 502 during the precomputation stage. Numerical tests have confirmed that this  
 503 error analysis is quite sharp. Constructive algorithms to find nearly optimal  
 504 parameters were proposed.

505 Since efficient FFT algorithms are available in virtually every computing en-  
 506 vironment, the time-critical stages of the algorithm are much easier to im-  
 507 plement. Furthermore, although the asymptotic arithmetic complexity is un-  
 508 changed, the smaller constant in the FFT's  $\mathcal{O}(N \log N)$  complexity yields  
 509 improved running times.

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## 578 A Asymptotic Complexity

579 In order to resolve a sufficient number of spherical harmonics, the number  
580 of points in the  $\phi$  and  $\theta$  directions must be  $\mathcal{O}(\ell) = \mathcal{O}(\kappa a)$ , where  $a$  is the  
581 side length of the box. Therefore, the total number of quadrature points is  
582  $\mathcal{O}(\ell^2) = \mathcal{O}((\kappa a)^2)$ . If  $a_0$  is the side length of a box at the root of the octree  
583 (level  $l = 0$ ), then the side length of a box at level  $l$  is  $a_l = 2^{-l}a_0$  and has  
584  $\mathcal{O}((\kappa 2^{-l}a_0)^2)$  quadrature points.

585 Suppose the octree has levels  $l = 0, \dots, L$  with boxes of side length  $a_l = 2^{-l}a_0$ .  
586 A box at level  $l$  will contain a field approximation defined over a quadrature

587 of size  $\mathcal{O}((\kappa 2^{-l} a_0)^2)$ . We now determine the number of operations required for  
 588 each step.

589 (1) **Initialization/Collection:** This step requires sampling  $e^{i\kappa \hat{s} \cdot \mathbf{r}}$  for each  
 590 of the  $N$  source points at the leaves of the tree. Thus, this step is

$$\mathcal{O}(N(\kappa 2^{-L} a_0)^2) = \mathcal{O}(N 2^{-2L} (\kappa a_0)^2)$$

591 (2) **Upward Pass:** This step requires aggregating and interpolating each  
 592 outgoing field upward through the tree. The type of interpolation algo-  
 593 rithm is key to the running time of this step.

594 At level  $l$  in the tree, the number of interpolations that must be per-  
 595 formed is equal to the number of boxes at that level in the tree. The  
 596 number of boxes depends on the distribution of source points. If the  
 597 source points are uniformly distributed over a volumetric domain then  
 598 the asymptotic number of boxes at level  $l$  in the octree is  $\mathcal{O}(8^l)$ . How-  
 599 ever, if the source points are uniformly distributed over the surface of an  
 600 object then the asymptotic number of boxes at level  $l$  is  $\mathcal{O}(4^l)$ .

601 • **Direct method:** Each direct interpolation requires  $\mathcal{O}(K_L K_{L-1}) =$   
 602  $\mathcal{O}((\kappa 2^{-l} a_0)^2 (\kappa 2^{-(l-1)} a_0)^2)$  operations. Thus, the direct method has com-  
 603 plexity

$$\begin{aligned} \text{Volume:} & \sum_{l=3}^L \mathcal{O}(8^l (\kappa 2^{-l} a_0)^2 (\kappa 2^{-(l-1)} a_0)^2) = \mathcal{O}((\kappa a_0)^4) \\ \text{Surface:} & \sum_{l=3}^L \mathcal{O}(4^l (\kappa 2^{-l} a_0)^2 (\kappa 2^{-(l-1)} a_0)^2) = \mathcal{O}((\kappa a_0)^4) \end{aligned}$$

604 • **Fast global methods:** By using a fast interpolation method, the com-  
 605 plexity for an individual interpolation is reduced to  $\mathcal{O}(K_l \log(K_l))$ . The  
 606 upward pass complexity then becomes

$$\begin{aligned} \text{Volume:} & \sum_{l=3}^L \mathcal{O}(8^l (\kappa 2^{-l} a_0)^2 \log(\kappa 2^{-l} a_0)) = \mathcal{O}(2^L (\kappa a_0)^2 (\log(\kappa a_0) - L)) \\ \text{Surface:} & \sum_{l=3}^L \mathcal{O}(4^l (\kappa 2^{-l} a_0)^2 \log(\kappa 2^{-l} a_0)) = \mathcal{O}(L (\kappa a_0)^2 (\log(\kappa a_0) - L)) \end{aligned}$$

607 • **Local methods:** Local methods use a stencil of some given size to  
 608 compute the interpolated values. Although these methods introduce  
 609 additional error, they have fast execution times with  $\mathcal{O}(K_l)$  operations.  
 610 The upward pass complexity then becomes

$$\begin{aligned} \text{Volume:} & \sum_{l=3}^L \mathcal{O}(8^l (\kappa 2^{-l} a_0)^2) = \mathcal{O}(2^L (\kappa a_0)^2) \\ \text{Surface:} & \sum_{l=3}^L \mathcal{O}(4^l (\kappa 2^{-l} a_0)^2) = \mathcal{O}(L (\kappa a_0)^2) \end{aligned}$$

611 (3) **Transfer Pass:** Since each box has a maximum of 189 transfers and the  
 612 transfer function is diagonal, the running time of this step is

$$\begin{aligned} \text{Volume:} & \quad \sum_{l=2}^L \mathcal{O}(8^l (\kappa 2^{-l} a_0)^2) = \mathcal{O}(2^L (\kappa a_0)^2) \\ \text{Surface:} & \quad \sum_{l=2}^L \mathcal{O}(4^l (\kappa 2^{-l} a_0)^2) = \mathcal{O}(L (\kappa a_0)^2) \end{aligned}$$

613 (4) **Downward Pass:** The downward pass is the adjoint operation of the  
 614 Upward Pass and has the same asymptotic complexity.

615 (5) **Finalization:** For each of the  $N$  field points, we integrate the spherical  
 616 function at the leaf and compute the close contributions from neighboring  
 617 boxes:

$$\mathcal{O}(N 2^{-2L} (\kappa a_0)^2) + \mathcal{O}(\text{close})$$

618 In the worst case, the close interaction is  $\mathcal{O}(N^2)$  which occurs when there  
 619 is an accumulation of points somewhere in the domain. In that case a  
 620 different scheme is required since the expansions used in this paper are  
 621 unstable at low frequency. When the field points are distributed roughly  
 622 uniformly, then

$$\begin{aligned} \text{Volume:} & \quad \mathcal{O}(\text{close}) = \mathcal{O}((N/8^L)^2) = \mathcal{O}(2^{-6L} N^2) \\ \text{Surface:} & \quad \mathcal{O}(\text{close}) = \mathcal{O}((N/4^L)^2) = \mathcal{O}(2^{-4L} N^2) \end{aligned}$$

623 The total asymptotic running time then depends on the scaling of  
 624 the number of points with  $a_0$ , the scaling of  $L$  with  $N$  or  $a_0$ , and the  
 625 interpolation methods that are used.

626 For a volume of scatters, we have  $\mathcal{O}(N) = \mathcal{O}((\kappa a_0)^3)$  and let  $L \sim$   
 627  $\log(N^{1/3})$  to achieve a total algorithmic complexity of  $\mathcal{O}(N \log N)$  using  
 628 fast global interpolation methods and  $\mathcal{O}(N)$  by using approximate, local  
 629 methods. For a surface of scatters, we have  $\mathcal{O}(N) = \mathcal{O}((\kappa a_0)^2)$  and let  
 630  $L \sim \log(N^{1/2})$  to achieve a total algorithmic complexity of  $\mathcal{O}(N \log^2 N)$   
 631 using fast global interpolation methods and  $\mathcal{O}(N \log N)$  by using approx-  
 632 imate, local methods.